

2SAT_v

$\neg$ 2SAT

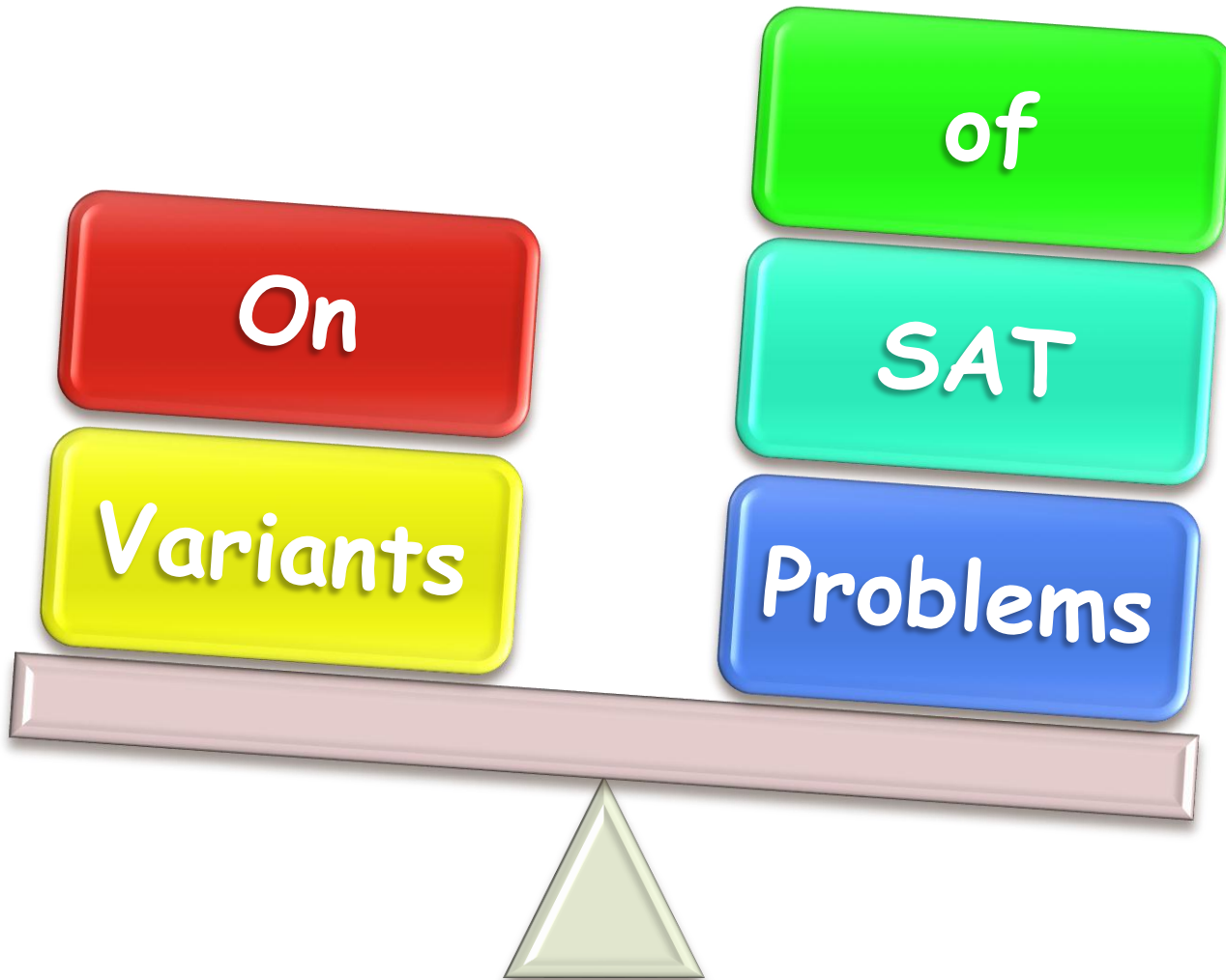
On

Variants

of

SAT

Problems



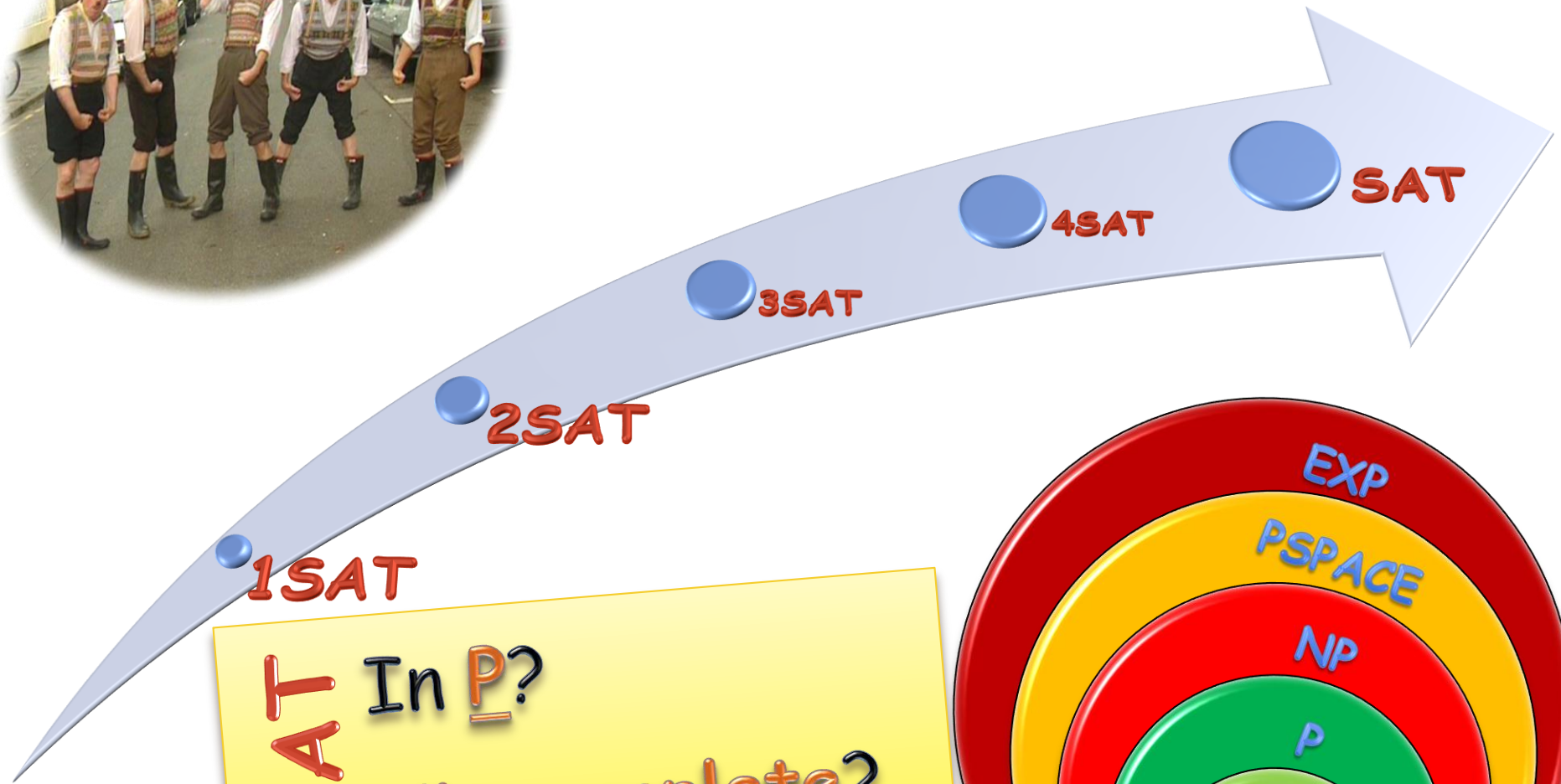
Goal:

- Discuss the complexity of variants of SAT

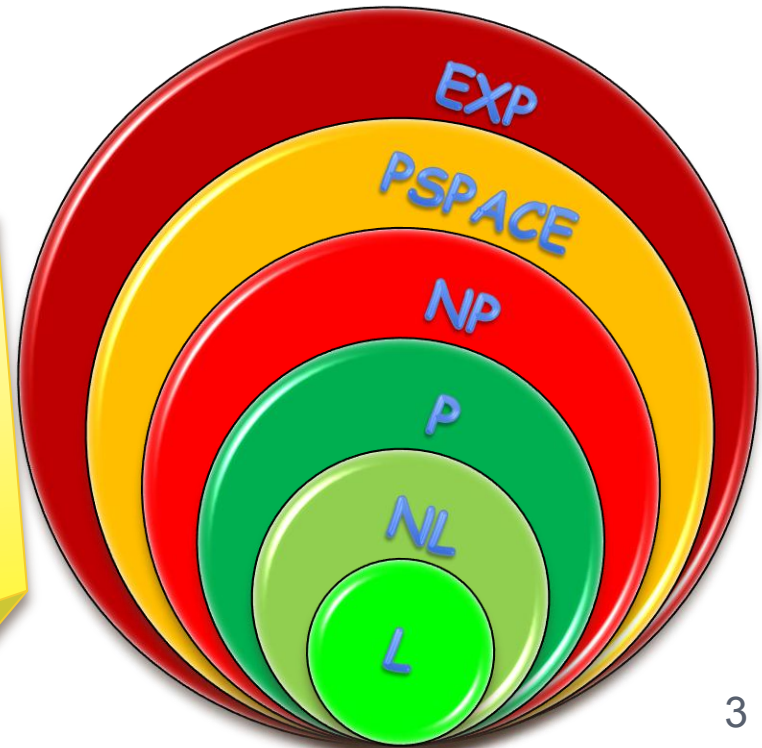
Plan:

- General
- 2SAT
- Max2SAT

Special cases of SAT



2SAT In P?
NL-complete?
Variants?



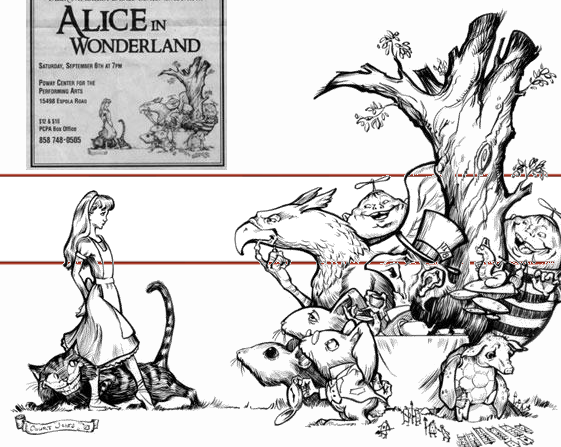
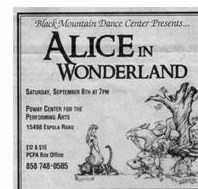
2SAT Instance:

- a 2-CNF formula ϕ

$$\phi \quad (\neg x \vee y) \wedge (\neg y \vee z) \wedge (x \vee \neg z)$$

Decision Problem:

- is ϕ satisfiable?

Theorem:

- $2SAT \in P$

Proof:

- Reduce 2SAT to a graph problem in P:
construct G_ϕ -- then specify problem

Implication graph $G_\varphi = (V_\varphi, E_\varphi)$

V_φ

- 1 vertex for every literal of φ

E_φ

note

edges: $(\alpha, \beta) \in E_\varphi \Leftrightarrow (\neg\beta, \neg\alpha) \in E_\varphi$
paths: $\alpha \mapsto \beta \Leftrightarrow \neg\beta \mapsto \neg\alpha$

- edge $(\alpha, \beta) \Leftrightarrow \varphi$ contains clause $(\neg\alpha \vee \beta)$

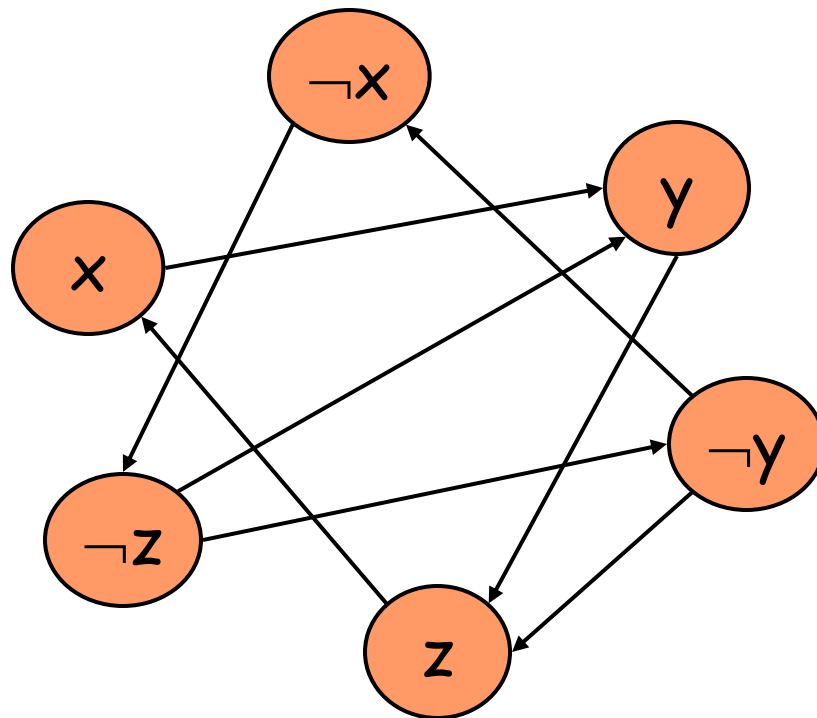
Theorem:

note $\alpha \mapsto \beta \Rightarrow$
 $\alpha \Rightarrow \beta$

- φ is unsatisfiable $\Leftrightarrow$
 $\exists x$ s.t. $x \mapsto \neg x$ and $\neg x \mapsto x$ in G_φ

Implication graph : Example

$$(\neg x \vee y) \wedge (\neg y \vee z) \wedge (x \vee \neg z) \wedge (z \vee y)$$



Correctness

Completeness:

- $x \mapsto \neg x \Rightarrow$ can't assign **TRUE** to x
- $\neg x \mapsto x \Rightarrow$ can't assign **FALSE** to x

Soundness:

- Repeat
Pick an x ; if $x \mapsto \neg x$, $\alpha = \neg x$ o/w $\alpha = x$ -
no $\alpha \mapsto \neg \alpha$, hence assign **TRUE** to α ;
Then, $\forall$ literal β s.t. $\alpha \mapsto \beta$:
assign **TRUE** to β and **FALSE** to $\neg \beta$
- No inconsistencies!

note

$$\alpha \mapsto \beta \wedge \alpha \mapsto \neg \beta \\ \Rightarrow \alpha \mapsto \neg \alpha$$

Graph Connectivity (CONN)

CONN Instance:

- a directed graph $G=(V,E)$ and 2 vertices $s,t \in V$

Decision Problem:

- Is there is a path from s to t in G ?

Theorem:

- $CONN \in P$

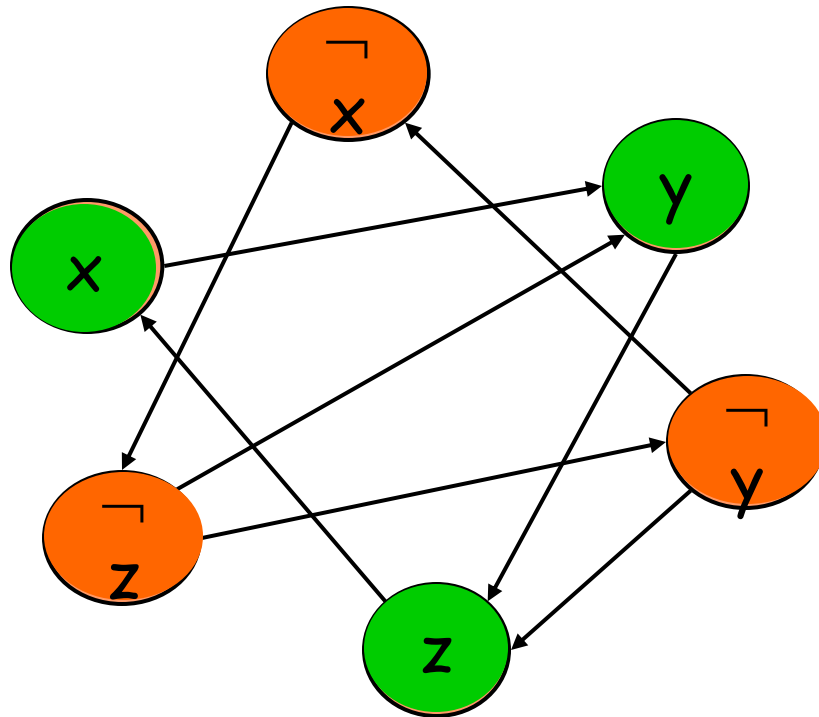
Apply some search algorithm (DFS/BFS)

Corollary:

- "Ex s.t. $x \mapsto \neg x$ and $\neg x \mapsto x$ in G_ϕ " $\in P$ ■

An Assignment: example

- Construct an assignment as follows:



Theorem:

- Max2SAT is NP-hard

note clearly
Max2SAT $\in$ NP

Proof: $3SAT \leq_p \text{Max2SAT}$

- Replace each $C = (\alpha \vee \beta \vee \gamma)$ of ϕ w/ 10 clauses in ϕ' :
 $(\alpha) \wedge (\beta) \wedge (\gamma) \wedge (w_C) \wedge (\neg \alpha \vee \neg \beta) \wedge (\neg \beta \vee \neg \gamma) \wedge (\neg \gamma \vee \neg \alpha) \wedge$
 $(\alpha \vee \neg w_C) \wedge (\beta \vee \neg w_C) \wedge (\gamma \vee \neg w_C).$
- Set $K = 7|\phi|$.

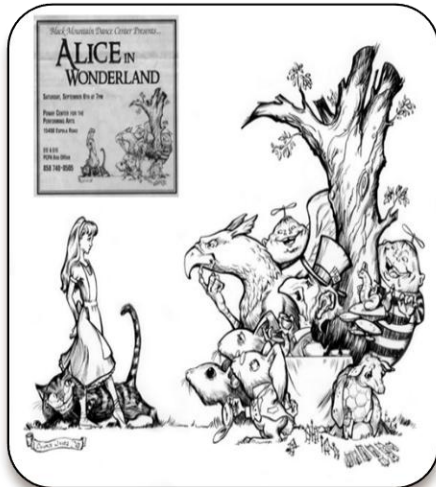
note $w_C = "a = b = c = \text{TRUE}?"$
maximizes satisfiab.

Completeness:

- $C = (\alpha \vee \beta \vee \gamma)$ satisfied $\Rightarrow$ 7/10 clauses satisfied

Soundness:

- $C = (\alpha \vee \beta \vee \gamma)$ unsatisfied $\Rightarrow \leq 6/10$ clauses satisfied



Discussed variants of **SAT**
Also: **Maximization** Problems



Special cases of **NPC** problems may
be in **P**: **SAT** vs. **2SAT**

Optimization versions of problems
in **P** may be hard: **2SAT** vs. **Max-2-SAT**

WWindex

SAT

Max-2-SAT

NPC



2SAT

Max-2-SAT

NPC



NL Complete

NP-Hard

Complexity
Classes

NP

NL

P

L

co-NP

EXPTIME

PSPACE



Papadimitriou,
Christos