

HIERARCHICAL MODEL OF THE INTERACTION OF MICROPROGRAMMED AUTOMATA

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Matematika i Vychislitel'naya Tekhnika,
21, No. 3, pp. 77-83, 1987

ISSN 0013-181.48:512.56

A model of interaction of microprogrammed automata at the nodes of a hierarchical structure is suggested. The model allows minimizing in the constituent automata the number of inputs and outputs for communication signals; it is designed for decompositional synthesis of automata based on PLA with internal memory.

In the synthesis of control devices based on PLA with a memory (PLAM), methods for decomposition of microprogrammed automata (MPA) are needed that are adjusted to the specifics of the system base, including the rigid constraints imposed by the external parameters (the number of inputs and outputs) on the parameters of the constituent automata. Decomposition methods that minimize the number of communication signals sent between constituent automata may be useful in this respect. One such method [1], introduces into the circuit a communication automaton which controls the constituent automata. This method takes into account the specifics of PLAM as the core unit, but does not help accommodate the specifics of the automata. Each of the constituent automata in a network is much more complex than a communication automaton as measured by the number of input terminals, transitions, and states; using a single board PLA for each communication automaton is therefore unjustified.

In this paper we propose a hierarchical model of interaction of constituent automata that overcomes these shortcomings.

The constituent automata in the model are the nodes of a hierarchical structure. A circuit with communication automata [1] is a two-layer hierarchical structure. In addition to a greater number of control levels, it differs from the model of [1] in that it allows the possibility for each constituent automaton of the network to realize the input and output functions of the initial automaton as well as the functions of the constituent automaton of the next hierarchical layer.

We will introduce definitions for the description of the hierarchical model of the interaction of automata.

We will say that the root automaton is the "boss" of the roots of its subtrees; the others are called "subordinates" of their "boss."

The relation of subordination on the set of constituent automata is a partial order relation. The fact that an automaton S^i is the "boss" of the automaton S^j - not necessarily the immediate boss - is written as $S^i > S^j$.

An automaton S^m of the hierarchical circuit is called the "immediate boss" (IB) of the automaton S^i , denoted $S^m = IB(S^i)$, if $S^m > S^i, S^n > S^i$ implies that either $S^n = S^m$ or $S^n = S^i$. The automaton S^i in that case is called the "immediate subordinate" (IS) of S^m . The set of automata that are immediate subordinates of S^m will be denoted by $IS(S^m)$.

$S^i \in IS(S^m)$ if and only if $S^m = IB(S^i)$.

The least upper bound [2] S^k for the automata S^i and S^j will be called the "nearest common boss" (NCB) of S^i and S^j : $S^k = NCB(S^i, S^j)$.

According to this terminology, each automaton of the hierarchical network is a boss

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Table 1

No.	a_i	a_j	$X(a_i, a_j)$	$Y(a_i, a_j)$
1	a_1	a_2	$\bar{x}_2 \bar{x}_4$	$y_1 y_6$
2		a_3	$x_1 \bar{x}_2$	$y_1 y_6$
3		a_3	$\bar{x}_1 \bar{x}_2 \bar{x}_4$	$y_6 y_6$
4		a_3	$x_2 \bar{x}_3 \bar{x}_4$	$y_6 y_6$
5		a_4	$x_2 x_3$	y_1
6		a_5	$x_2 \bar{x}_3 \bar{x}_4$	$y_1 y_6$
7	a_2	a_6	$x_1 x_3$	y_1
8		a_7	\bar{x}_3	$y_1 y_6$
9		a_7	\bar{x}_1	$y_1 y_6$
10	a_3	a_4	1	y_6
11	a_4	a_6	$\bar{x}_6 \bar{x}_7 \bar{x}_8$	$y_6 y_6$
12		a_6	$x_6 x_8$	y_1
13		a_7	$\bar{x}_6 \bar{x}_8$	$y_1 y_6$
14		a_8	$x_6 \bar{x}_8$	$y_1 y_6 y_6$
15		a_8	$\bar{x}_6 \bar{x}_7 \bar{x}_8$	$y_1 y_6 y_6$
16	a_5	a_5	$\bar{x}_7 \bar{x}_9 \bar{x}_3$	$y_6 y_7 y_6$
17		a_7	$x_7 x_9$	y_6
18		a_8	$\bar{x}_7 \bar{x}_9$	$y_7 y_6$
19		a_9	$x_7 \bar{x}_9$	—
20		a_{10}	$\bar{x}_2 \bar{x}_3 \bar{x}_7$	y_7
21	a_4	a_2	$x_3 x_4 \bar{x}_8$	y_6
22		a_3	$\bar{x}_3 \bar{x}_4$	$y_1 y_3$
23		a_4	$x_3 x_4 x_5 \bar{x}_6 \bar{x}_7 \bar{x}_8$	$y_2 y_4$
24		a_7	$\bar{x}_4 \bar{x}_5 \bar{x}_7$	y_1
25		a_9	$\bar{x}_4 \bar{x}_5 \bar{x}_6$	$y_1 y_4$
26		a_{10}	$\bar{x}_4 \bar{x}_5 \bar{x}_6$	$y_2 y_4$
27		a_{10}	$\bar{x}_4 \bar{x}_6 \bar{x}_7$	$y_2 y_4$
28		a_{11}	$\bar{x}_4 \bar{x}_5 \bar{x}_6 \bar{x}_7$	$y_1 y_2 y_3$
29		a_{11}	$x_3 x_4 x_5 \bar{x}_6$	$y_1 y_2 y_3$
30		a_{11}	$x_3 x_4 \bar{x}_7 \bar{x}_8$	$y_1 y_2 y_3$
31		a_{11}	$x_3 x_4 x_5 \bar{x}_6$	$y_1 y_2 y_3$
32	a_7	a_8	1	y_1
33	a_8	a_2	\bar{x}_1	—
34		a_{11}	x_1	y_6
35	a_9	a_{13}	x_2	y_1
36		a_{14}	\bar{x}_2	y_5
37	a_{10}	a_{15}	x_1	y_6
38		a_{16}	\bar{x}_1	$y_5 y_6$
39	a_{11}	a_1	\bar{x}_9	y_6
40		a_2	$x_9 \bar{x}_8$	$y_7 y_6$
41		a_{17}	$x_9 \bar{x}_8$	$y_6 y_7$
42	a_{12}	a_3	$x_2 x_3 \bar{x}_7$	$y_7 y_6$
43		a_4	$x_2 x_3 \bar{x}_7$	y_6
44		a_5	$x_2 \bar{x}_3$	y_7
45		a_6	\bar{x}_2	$y_7 y_6$
46	a_{13}	a_7	$\bar{x}_3 \bar{x}_6 \bar{x}_8$	$y_6 y_{10}$
47		a_7	$x_6 \bar{x}_8$	$y_6 y_{10}$
48		a_8	$\bar{x}_4 \bar{x}_9$	$y_6 y_{10}$
49		a_9	$x_4 \bar{x}_9$	$y_6 y_6 y_{10}$
50		a_9	$\bar{x}_4 \bar{x}_5 \bar{x}_6$	$y_6 y_6 y_{10}$
51	a_{14}	a_1	1	y_7
52	a_{15}	a_1	$x_7 x_8$	$y_6 y_6$
53		a_4	$x_7 \bar{x}_8$	$y_7 y_6$
54		a_7	\bar{x}_7	y_6
55	a_{16}	a_{11}	1	y_5
56	a_{17}	a_{12}	1	$y_6 y_7 y_6$

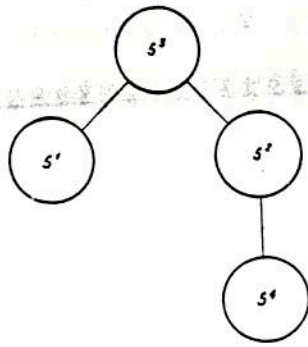


Fig. 1. Network structure in the decomposition of MPA S.

Table 2
Constituent Automaton S^1

a_i	a_j	$X(a_i, a_j)$	$Q_j(S^1)$	$Y(a_i, a_j)$	$P_j(S^1)$
a_4	b_1	$x_2 x_4 \bar{x}_8$	—	y_0	$P_2(S^1) = q_{31} q_{22} q_{33}$
	b_2	$\bar{x}_3 x_4$	—	$y_1 y_3$	$P_3(S^1) = q_{31} q_{22} \bar{q}_{33}$
	a_4	$x_2 x_4 x_5 \bar{x}_6 x_7 x_8$	—	$y_2 y_4$	$P_4(S^1) = \bar{q}_{31} \bar{q}_{22} \bar{q}_{33}$
	b_1	$\bar{x}_4 \bar{x}_5 x_7$	—	y_1	$P_7(S^1) = q_{31} \bar{q}_{22} q_{33}$
	b_1	$\bar{x}_4 x_5 x_7$	—	$y_1 y_4$	$P_8(S^1) = q_{31} \bar{q}_{22} \bar{q}_{33}$
	b_1	$\bar{x}_4 x_5 x_6$	—	$y_2 y_4$	$P_{10}(S^1) = \bar{q}_{31} q_{22} q_{33}$
	b_1	$\bar{x}_4 \bar{x}_6 x_7$	—	$y_2 y_4$	$P_{10}(S^1) = \bar{q}_{31} q_{22} \bar{q}_{33}$
	b_1	$\bar{x}_4 \bar{x}_6 x_7$	—	$y_1 y_2 y_3$	$P_{11}(S^1) = \bar{q}_{31} q_{22} \bar{q}_{33}$
	b_1	$\bar{x}_4 \bar{x}_6 x_7$	—	$y_1 y_2 y_3$	$P_{11}(S^1) = \bar{q}_{31} q_{22} \bar{q}_{33}$
	b_1	$x_2 x_4 \bar{x}_5 x_8$	—	$y_1 y_2 y_3$	$P_{11}(S^1) = \bar{q}_{31} q_{22} \bar{q}_{33}$
	b_1	$x_2 x_4 \bar{x}_5 x_8$	—	$y_1 y_2 y_3$	$P_{11}(S^1) = \bar{q}_{31} q_{22} \bar{q}_{33}$
	b_1	$x_2 x_4 \bar{x}_5 x_8$	—	$y_1 y_2 y_3$	$P_{11}(S^1) = \bar{q}_{31} q_{22} \bar{q}_{33}$
b_1	a_4	—	$Q_0(S^1) = q_{11}$	—	—
	b_1	—	$Q_0(S^1) = \bar{q}_{11}$	—	—

subordinate at the same time. The exceptions are the automata of the last level, which are all subordinates, and the automaton of the first level, which is only a boss.

The hierarchical network operates as follows.

During the operation of a particular constituent automaton, all the automata in the network are in the special state of "waiting for a response" from particular subordinate automata. The other automata of the network are in the state of "waiting for a response" from the boss automaton.

Each constituent automaton except for the root automaton and the leaf automata, is connected with its boss and its subordinates. When transferring control to a subordinate, the constituent automaton goes into the state of waiting for its response. When transferring control to a boss the automaton goes into the state of response waiting.

We can define the constituent automata of a hierarchical network in formal terms. The description will be illustrated by the decomposition of MPA $S(A, X, Y, \delta, \lambda, a_1)$ (A - is the set of states, X - are the sets of input and output signals, δ, λ - are the transition functions, respectively, and a_1 - is the initial state); it is defined by Table 1.

Suppose that on the set A of the states of MPA S a partition $\pi = \{A^1, \dots, A^v\}$ is given (v is the number of partition blocks). In the example $\pi = \{A^1, A^2, A^3, A^4\} = \{6, 1, 2, 3, 7, 8, 9, 10, 11, 12, 13, 14, 15, 16, 17, 4, 13\}$. We will perform the decomposition of this MPA according to the structure shown in Fig. 1.

The set B^u of the states of a constituent MPA S^u consists of the set A^u of the states of the corresponding partition block, the state b_u of request waiting, and the set of the states of response waiting: $B^u = A^u \cup \{b_u\} \cup C^u$, where $C^u = \bigcup_{S^i \in \pi(S^u)} c_u(S^i)$, $c_u(S^i)$ is the state of response waiting of the subordinate MPA S^i . In this example: $B^1 = \{a_0, b_1\}$; $B^2 = \{a_5, a_{11}, a_{12}, a_{14}, a_{15}, a_{16}, a_{17}, c_2(S^1), c_3(S^2)\}$; $B^3 = \{a_4, a_{13}, b_4\}$; $B^4 = \{a_3, a_7, a_8, a_9, a_{10}, b_2, c_2(S^4)\}$.

Table 3
Constituent MPA S^2

a_i	a_j	$X(a_i, a_j)$	$\hat{Q}_j(S^2)$	$Y(a_i, a_j)$	$\hat{P}_j(S^2)$
a_1	a_2	\bar{x}_2, \bar{x}_4	—	$y_1 y_5$	—
	a_3	\bar{x}_1, \bar{x}_2	—	$y_1 y_5$	—
	a_5	$\bar{x}_1, \bar{x}_2, \bar{x}_4$	—	$y_2 y_6$	—
	a_7	$\bar{x}_2, \bar{x}_3, \bar{x}_4$	—	$y_2 y_6$	—
	$c_2(S^1)$	\bar{x}_2, \bar{x}_3	—	y_1	$P_1(S^2) = q_{11} q_{42}$
	b_2	$\bar{x}_2, \bar{x}_3, \bar{x}_4$	—	$y_1 y_6$	$P_2(S^2) = q_{21} q_{22} q_{33}$
a_2	b_2	\bar{x}_1, \bar{x}_3	—	y_1	$P_8(S^2) = q_{21} q_{22} q_{33}$
	a_7	\bar{x}_3	—	$y_1 y_6$	—
	a_7	\bar{x}_1	—	$y_1 y_6$	—
a_6	a_6	\bar{x}_1	—	—	—
	b_2	\bar{x}_1	—	y_6	$P_{11}(S^2) = q_{31} q_{32} q_{33}$
a_9	$c_2(S^1)$	\bar{x}_2	—	y_1	$P_{13}(S^2) = q_{41} q_{42}$
	b_2	\bar{x}_2	—	y_6	$P_{14}(S^2) = q_{31} q_{32} q_{33}$
a_{10}	b_2	\bar{x}_1	—	y_6	$P_{15}(S^2) = q_{31} q_{32} q_{33}$
	b_2	\bar{x}_1	—	$y_2 y_6$	$P_{16}(S^2) = q_{31} q_{32} q_{33}$
a_7	a_9	1	—	y_1	—
	$c_2(S^1)$	1	—	y_6	$P_4(S^2) = q_{41} q_{42}$
b_2	a_1	—	$Q_1(S^2) = q_{21} q_{22} q_{23} q_{24}$	—	—
	a_2	—	$Q_2(S^2) = q_{21} q_{22} q_{23} q_{24}$	—	—
	a_6	—	$Q_8(S^2) = q_{21} q_{22} q_{23} q_{24}$	—	—
	a_9	—	$Q_9(S^2) = q_{21} q_{22} q_{23} q_{24}$	—	—
	a_{10}	—	$Q_{10}(S^2) = q_{21} q_{22} q_{23} q_{24}$	—	—
	a_7	—	$Q_7(S^2) = q_{21} q_{22} q_{23} q_{24}$	—	—
	a_3	—	$Q_3(S^2) = q_{21} q_{22} q_{23} q_{24}$	—	—
	a_5	—	$Q_5(S^2) = q_{21} q_{22} q_{23} q_{24}$	—	—
	$c_2(S^1)$	—	$Q_{13}(S^2) = q_{21} q_{22} q_{23} q_{24}$	—	$P_{10}(S^2) = q_{41} q_{42}$
	$c_2(S^1)$	—	$Q_4(S^2) = q_{21} q_{22} q_{23} q_{24}$	—	$P_4(S^2) = q_{41} q_{42}$
	$c_2(S^1)$	—	$Q_6(S^2) = q_{21} q_{22} q_{23} q_{24}$	—	—
	$c_2(S^1)$	—	$Q_8(S^2) = q_{21} q_{22} q_{23} q_{24}$	—	—
$c_2(S^1)$	a_7	—	$q_{21} q_{22} q_{23}$	—	—
	a_6	—	$q_{21} q_{22} q_{23}$	—	—
	a_9	—	$q_{21} q_{22} q_{23}$	—	—
	b_2	—	$q_{21} q_{22} q_{23}$	—	$P_5(S^2) = q_{31} q_{32} q_{33}$
	b_2	—	$q_{21} q_{22} q_{23}$	—	$P_6(S^2) = q_{31} q_{32} q_{33}$
	$c_2(S^1)$	—	$q_{21} q_{22} q_{23}$	—	—

2. Put into correspondence to each block A^u of the partition π a set of input variables \tilde{X}^u interrogated at the transitions from the states of this block: $\tilde{X}^u = \bigcup_{a_i \in A^u} X(a_i)$.

The input variables of the constituent MPA S^u are the elements of the set X^u and the elements of the set $Q(S^u) = \{q_{u1}, \dots, q_{un}\}$ of the additional variables sent to the input S^u : $X^u = \tilde{X}^u \cup Q(S^u)$. In the example: $X^1 = X^1 \cup Q(S^1) = \{x_3, x_4, x_5, x_6, x_7, x_8, q_{11}\}$, $X^2 = \tilde{X}^2 \cup Q(S^2) = \{x_1, x_2, x_3, x_4, q_{21}, q_{22}, q_{23}, q_{24}\}$, $X^3 = \{x_2, x_3, x_7, x_8, q_{31}, q_{32}, q_{33}\}$, $X^4 = \{x_3, x_4, x_5, x_6, x_7, q_{41}, q_{42}\}$.

3. For each block A^u of the partition π , we form a set \tilde{Y}^u of the output variables generated at the transitions from the states of this block: $\tilde{Y}^u = \bigcup_{a_i \in A^u} Y(a_i)$. The output variables of the constituent MPA S^u are the elements of the set \tilde{Y}^u and the elements of the set $P(S^u)$ of the additional variables sent to the inputs of the automaton IB (S^u) and each of IS (S^u):

$$Y^u = \tilde{Y}^u \cup P(S^u); P(S^u) = \bigcup_{S^i \in IS(S^u)} Q(S^i) \cup Q(IB(S^u)).$$

In the example:

$$Y^1 = \tilde{Y}^1 \cup Q(S^3); Y^2 = \tilde{Y}^2 \cup Q(S^3) \cup Q(S^4);$$

$$Y^3 = \tilde{Y}^3 \cup Q(S^1) \cup Q(S^2); Y^4 = \tilde{Y}^4 \cup Q(S^2).$$

4. We define the functions of the transitions δ^u and the outputs λ^u of the constituent automaton S^u . Suppose that in the initial MPA S there is a transition from a_1 to a_2 .

Table 4

Constituent MPA S³

a_i	a_j	$X(a_i, a_j)$	$\hat{Q}_j(S^i)$	$Y(a_i, a_j)$	$\hat{P}_j(S^i)$
a_{15}	$c_3(S^2)$	x_7x_8	—	y_6y_8	$q_{21}q_{22}q_{23}q_{24}$
	$c_3(S^1)$	$x_7\bar{x}_8$	—	y_7y_8	q_{11}
	$c_3(S^2)$	\bar{x}_7	—	y_6	—
a_5	$c_3(S^1)$	$\bar{x}_7\bar{x}_2x_3$	—	$y_6y_7y_8$	q_{11}
	$c_3(S^2)$	x_7x_2	—	y_6	$q_{21}q_{22}q_{23}q_{24}$
	$c_3(S^2)$	\bar{x}_7x_2	—	y_7y_8	$q_{21}q_{22}q_{23}q_{24}$
	$c_3(S^2)$	$x_7\bar{x}_2$	—	—	$q_{21}q_{22}q_{23}q_{24}$
	$c_3(S^2)$	$\bar{x}_7\bar{x}_2\bar{x}_3$	—	y_7	$q_{21}q_{22}q_{23}q_{24}$
a_{12}	$c_3(S^2)$	$x_2x_3x_7$	—	y_7y_8	$q_{21}q_{22}q_{23}q_{24}$
	$c_3(S^2)$	$x_2x_3\bar{x}_7$	—	y_8	$q_{21}q_{22}q_{23}q_{24}$
	a_6	$x_2\bar{x}_3$	—	y_7	—
	$c_3(S^1)$	\bar{x}_3	—	y_7y_8	q_{11}
a_{11}	$c_3(S^2)$	\bar{x}_2	—	y_6	$q_{21}q_{22}q_{23}q_{24}$
	$c_3(S^2)$	$x_2\bar{x}_8$	—	y_7y_8	$q_{21}q_{22}q_{23}q_{24}$
	a_{17}	x_2x_8	—	y_6y_7	—
a_{14}	$c_3(S^2)$	1	—	y_7	$q_{21}q_{22}q_{23}q_{24}$
a_{16}	a_{11}	1	—	y_8	—
a_{17}	a_{12}	1	—	$y_6y_7y_8$	—
$c_3(S^1)$	a_{11}	—	$q_{31}q_{32}q_{33}$	—	—
	$c_3(S^2)$	—	$q_{31}q_{32}q_{33}$	—	$q_{21}q_{22}q_{23}q_{24}$
	$c_3(S^2)$	—	$q_{31}q_{32}q_{33}$	—	$q_{21}q_{22}q_{23}q_{24}$
	$c_3(S^2)$	—	$q_{31}q_{32}q_{33}$	—	$q_{21}q_{22}q_{23}q_{24}$
	$c_3(S^2)$	—	$q_{31}q_{32}q_{33}$	—	$q_{12}q_{22}q_{23}q_{24}$
	$c_3(S^2)$	—	$q_{31}q_{32}q_{33}$	—	$q_{21}q_{22}q_{23}q_{24}$
$c_3(S^2)$	a_{15}	—	$q_{31}q_{32}q_{33}$	—	—
	a_6	—	$q_{31}q_{32}q_{33}$	—	—
	a_{12}	—	$q_{31}q_{32}q_{33}$	—	—
	a_{11}	—	$q_{31}q_{32}q_{33}$	—	—
	a_{14}	—	$q_{31}q_{32}q_{33}$	—	—
	a_{16}	—	$q_{31}q_{32}q_{33}$	—	—
	c_1	—	$q_{31}q_{32}q_{33}$	—	q_{11}
	c_2	—	$q_{31}q_{32}q_{33}$	—	—

by the input signal X_n that results in the output signal Y_t : $\delta(a_i, X_n) = a_j$; $\lambda(a_i, X_n) = Y_t$, suppose that $a_i \in A^u$. Four alternatives are possible:

$a_j \in A^u$, then $\delta^u(a_i, X_n) = \delta(a_i, X_n) = a_j$; $\lambda^u(a_i, X_n) = \lambda(a_i, X_n) = Y_t$;

$a_j \in A^k$ ($k \neq u$), $S^k > S^u$, then in the automaton S^u : $\delta^u(a_i, X_n) = b_u$; $\lambda^u(a_i, X_n) = Y_t U P_j(S^u)$, where $P_j(S^u)$ is the function of additional input variables causing S^k to go to the state a_j . In all S^m ($S^k > S^m > S^u$): $\delta^m(c_m(S^u), Q(S^m)) = b_m$, $\lambda^m(c_m(S^u), Q(S^m)) = P_j(S^m)$;

$a_j \in A^k$ ($k \neq u$), $S^k < S^u$. Then in the automaton S^k : $\delta^k(a_i, X_n) = c_u(S^k)$, $\lambda^k(a_i, X_n) = Y_t U P_j(S^k)$. In the automaton S^k : $\delta^k(b_k, Q_j(S^k)) = \delta(a_i, X_n) = a_j$, $\lambda^k(b_k, Q_j(S^k)) = Y_0$. In all S^m ($S^k < S^m < S^u$): $\delta^m(b_m, Q(S^m)) = c_m(S^k)$, $\lambda^m(b_m, Q(S^m)) = P_j(S^m)$.

$a_j \in A^k$ ($k \neq u$), S^k and S^u do not lie on a common path to the root of the hierarchical tree. Let $S^r = \text{NIB}(S^k, S^u)$. Then, in the automaton S^u : $\delta^u(a_i, X_n) = b_u$, $\lambda^u(a_i, X_n) = Y_t U P_j(S^u)$. In the automaton S^k : $\delta^k(b_k, Q_j(S^k)) = a_j$, $\lambda^k(b_k, Q_j(S^k)) = Y_0$. In the automaton S^r : $\delta^r(c_r(S^u), Q_j(S^r)) = c_r(S^k)$, $\lambda^r(c_r(S^u), Q_j(S^r)) = P_j(S^u)$. In all S^m ($S^r > S^m > S^k$): $\delta^m(b_m, Q_j(S^m)) = c_m(S^k)$, $\lambda^m(b_m, Q_j(S^m)) = P_j(S^m)$. In all S^m ($S^r > S^m > S^u$): $\delta^m(c_m(S^u), Q_j(S^m)) = b_m$, $\lambda^m(c_m(S^u), Q_j(S^m)) = P_j(S^u)$.

Besides $\delta^u(b_u, Q_0^u) = b_u$, $\delta^u(c_u(S^k), Q_0^u) = c_u(S^k)$, $\lambda^u(b_u, Q_0^u) = \lambda^u(c_u(S^k), Q_0^u) = Y_0$, where Q_0^u is the conjunction of the variables of the set U^u , all of whose terms are inverted.

The number of additional input variables of the MPA S^u is defined as follows:

to the number of outputs due to an increased number of communication signals.

The interaction model always produces implementations with the number of PLAM boards at least not greater than model of [1], which is a particular case of the model suggested in this paper.

The use of a hierarchical interaction model in MPA decomposition made it possible to achieve the following compared with model of [1]: 1) reduce the number of interacting automata in the network; 2) take into account the specifics of the connections of the automata and reduce the number of communication variables by introducing a sequence in the connection between network automata; and 3) expand the class of circuits implementing control automata on a PLAM basis.

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